

PSEUDOALGEBRAS FOR THE PRESHEAF RELATIVE PSEUDOMONAD ARE COCOMPLETE

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1. RELATIVE PSEUDOMONADS AND THEIR PSEUDOALGEBRAS

Definition 1.1. Let \mathbb{C}, \mathbb{D} be bicategories and $J : \mathbb{D} \rightarrow \mathbb{C}$ a pseudofunctor. A *relative pseudomonad* $(T, i, *, \eta, \mu, \theta)$ along J comprises:

- for every object $X \in \mathbb{D}$ an object $TX \in \mathbb{D}$ and map $i_X : JX \rightarrow TX$ in \mathbb{C} , and
- a family of functors $(-)^*_{X,Y} : \mathbb{C}(JX, TY) \rightarrow \mathbb{C}(TX, TY)$ for $X, Y \in \mathbb{D}$,

along with three families of invertible 2-cells:

- $\eta_f : f \rightarrow f^*i$ for $f : JX \rightarrow TY$,
- $\mu_{f,g} : (f^*g)^* \rightarrow f^*g^*$ for $f : JX \rightarrow TY, g : JW \rightarrow TX$, and
- $\theta_X : i_X^* \rightarrow 1_{TX}$ for $X \in \mathbb{D}$,

such that the following two coherence diagrams commute:

(i) for $f : JX \rightarrow TY, g : JW \rightarrow TX, h : JV \rightarrow TW$,

$$(1) \quad \begin{array}{ccc} ((f^*g)^*h)^* & \xrightarrow{\mu_{f^*g,h}} & (f^*g)^*h^* \\ (\mu_{f,g,h})^* \downarrow & & \downarrow \mu_{f,g,h^*} \\ ((f^*g^*)h)^* & & (f^*g^*)h^* \\ \sim \downarrow & & \downarrow \sim \\ (f^*(g^*h))^* & \xrightarrow{\mu_{f,g^*h}} f^*(g^*h)^* \xrightarrow{f^*\mu_{g,h}} & f^*(g^*h^*) \end{array}$$

(ii) for $f : JX \rightarrow TY$,

$$(2) \quad \begin{array}{ccc} f^* & \xrightarrow{(\eta_f)^*} (f^*i)^* \xrightarrow{\mu_{f,i}} f^*i^* & \\ & \searrow \sim & \downarrow f^*\theta \\ & & f^*1 \end{array}$$

Definition 1.2. Let T be a relative pseudomonad along $J : \mathbb{D} \rightarrow \mathbb{C}$. A *pseudoalgebra* $(A, a; \tilde{a}, \hat{a})$ comprises:

- an object $A \in \mathbb{C}$,
- a family of functors $(-)^a_X : \mathbb{C}(JX, A) \rightarrow \mathbb{C}(TX, A)$ for $X \in \mathbb{D}$,

along with two families of invertible 2-cells

- $\tilde{a}_f : f \rightarrow f^a i$ for $f : JX \rightarrow A$,
- $\hat{a}_{f,g} : (f^a g)^a \rightarrow f^a g^*$ for $f : JX \rightarrow A, g : JW \rightarrow TX$,

such that the following two coherence diagrams commute:

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(i) for $f : JX \rightarrow A$, $g : JW \rightarrow TX$, $h : JV \rightarrow TW$,

$$(3) \quad \begin{array}{ccc} ((f^a g)^a h)^a & \xrightarrow{\hat{a}_{f^a g, h}} & (f^a g)^a h^* \\ (\hat{a}_{f, g h})^a \downarrow & & \downarrow \hat{a}_{f, g h^*} \\ ((f^a g^*) h)^a & & (f^a g^*) h^* \\ \sim \downarrow & & \downarrow \sim \\ (f^a (g^* h))^a & \xrightarrow{\hat{a}_{f, g^* h}} f^a (g^* h)^* \xrightarrow{f^a \mu_{g, h}} & f^a (g^* h^*) \end{array}$$

(ii) for $f : JX \rightarrow A$,

$$(4) \quad \begin{array}{ccc} f^a & \xrightarrow{(\tilde{a}_f)^a} (f^a i)^a & \xrightarrow{\hat{a}_{f, i}} f^a i^* \\ & \searrow \sim & \downarrow f^a \theta \\ & & f^a 1 \end{array}$$

Lemma 1.3. *Let \mathbb{A} be a pseudoalgebra over the relative pseudomonad T . Then for $f : JX \rightarrow A$, $g : JW \rightarrow TX$, the diagram*

$$(5) \quad \begin{array}{ccc} f^a g & \xrightarrow{\tilde{a}} (f^a g)^a i & \xrightarrow{\hat{a}} (f^a g^*) i \\ & \searrow f^a \eta & \downarrow \sim \\ & & f^a (g^* i) \end{array}$$

also commutes.

Proof. Since \tilde{a} is a natural isomorphism, the required coherence is equivalent to

$$\begin{array}{ccc} (f^a g)^a i & \xrightarrow{\tilde{a}} ((f^a g)^a i)^a i & \xrightarrow{\hat{a}} ((f^a g^*) i)^a i \\ & \searrow \eta & \downarrow \sim \\ & & (f^a (g^* i))^a i \end{array}$$

and so it suffices to show

$$\begin{array}{ccc} (f^a g)^a & \xrightarrow{\tilde{a}} ((f^a g)^a i)^a & \xrightarrow{\hat{a}} ((f^a g^*) i)^a \\ & \searrow \eta & \downarrow \sim \\ & & (f^a (g^* i))^a \end{array}$$

commutes. Using this instance of equation (3)

$$\begin{array}{ccccc} ((f^a g)^a i)^a & \xrightarrow{\hat{a}} & ((f^a g^*) i)^a & \xrightarrow{\sim} & (f^a (g^* i))^a \\ \downarrow \hat{a} & & & & \downarrow \hat{a} \\ (f^a g)^a i^* & \xrightarrow{\hat{a}} & (f^a g^*) i^* & \xrightarrow{\sim} & f^a (g^* i)^* \\ & & & & \downarrow \mu \\ (f^a g)^a i^* & \xrightarrow{\hat{a}} & (f^a g^*) i^* & \xrightarrow{\sim} & f^a (g^* i^*) \end{array}$$

the clockwise composite becomes

$$\begin{array}{c}
(f^a g)^a \xrightarrow{\hat{a}} ((f^a g)^a i)^a \xrightarrow{\hat{a}} (f^a g)^a i^* \xrightarrow{\hat{a}} (f^a g^*) i^* \\
\downarrow \sim \\
f^a (g^* i^*) \\
\downarrow \mu^{-1} \\
f^a (g^* i)^* \\
\downarrow \hat{a}^{-1} \\
(f^a (g^* i))^a
\end{array}$$

and now using equation (4) we can replace the first two maps in the clockwise composite to obtain:

$$\begin{array}{c}
(f^a g)^a \xrightarrow{\sim} (f^a g)^a 1 \xrightarrow{\theta^{-1}} (f^a g)^a i^* \xrightarrow{\hat{a}} (f^a g^*) i^* \\
\downarrow \sim \\
f^a (g^* i^*) \\
\downarrow \mu^{-1} \\
f^a (g^* i)^* \\
\downarrow \hat{a}^{-1} \\
(f^a (g^* i))^a
\end{array}$$

which we fill in with four naturality squares, a bicategory coherences and relative pseudomonad coherence (2):

$$\begin{array}{ccccc}
(f^a g)^a & \xrightarrow{\sim} & (f^a g)^a 1 & \xrightarrow{\theta^{-1}} & (f^a g)^a i^* \\
\hat{a} \downarrow & & \hat{a} \downarrow & & \downarrow \hat{a} \\
f^a g^* & \xrightarrow{\sim} & (f^a g^*) 1 & \xrightarrow{\theta^{-1}} & (f^a g^*) i^* \\
\parallel & & \downarrow \sim & & \downarrow \sim \\
f^a g^* & \xrightarrow{\sim} & f^a (g^* 1) & \xrightarrow{\theta^{-1}} & f^a (g^* i^*) \\
& & \downarrow \sim & & \downarrow \mu^{-1} \\
& & f^a g^* & \xrightarrow{\eta} & f^a (g^* i)^* \\
& & \hat{a}^{-1} \downarrow & & \downarrow \hat{a}^{-1} \\
& & (f^a g)^a & \xrightarrow{\eta} & (f^a (g^* i))^a
\end{array}$$

Here the anticlockwise composite is equal to $(f^a \eta_g)^a$, as desired. \square

2. THE PRESHEAF RELATIVE PSEUDOMONAD

Write Cat for the 2-category of small categories, and write CAT for the 2-category of locally-small categories. Since the category of presheaves on a small

category is in general only locally small, it is natural to ask whether the presheaf construction

$$X \mapsto \text{Psh } X := [X^{op}, \text{Set}]$$

can be given the structure of a relative pseudomonad along the inclusion 2-functor $J : \text{Cat} \rightarrow \text{CAT}$.

This is shown in [FGHW18] via the construction of a relative pseudoadjunction; the structure of a relative pseudomonad is given to Psh as follows:

- for an object $X \in \text{Cat}$ we have $\text{Psh } X \in \text{CAT}$ and unit map $y_X : X \rightarrow \text{Psh } X$ given by the Yoneda embedding,
- for $X, Y \in \text{Cat}$ and a functor $f : X \rightarrow \text{Psh } Y$, the extension $f^* : \text{Psh } X \rightarrow \text{Psh } Y$ is given by the left Kan extension of f along the Yoneda embedding

$$\begin{array}{ccc} X & \xrightarrow{y} & \text{Psh } X \\ & \searrow f & \downarrow \eta_f \\ & & \text{Psh } Y \end{array} \quad \begin{array}{c} \xrightarrow{\eta_f} \\ \downarrow f^* := \text{Lan}_y f \end{array}$$

which also defines the 2-cells $\eta_f : f \rightarrow f^*y$ (note that since the Yoneda embedding is fully faithful the maps η_f are invertible, as required),

- for $f : JX \rightarrow TY$ and $g : JW \rightarrow TX$, the 2-cell $\mu_{f,g} : (f^*g)^* \rightarrow f^*g^*$ is uniquely determined by the universal property of the left Kan extension:

$$\begin{array}{ccc} W & \xrightarrow{y} & \text{Psh } W \\ & \searrow g & \downarrow \eta_g \\ & & \text{Psh } X \\ & & \downarrow f^* \\ & & \text{Psh } Y \end{array} \quad \begin{array}{ccc} W & \xrightarrow{y} & \text{Psh } W \\ & \searrow f^*g & \downarrow \eta_{f^*g} \\ & & (f^*g)^* \\ & & \downarrow \mu_{f,g} \\ & & \text{Psh } X \\ & & \downarrow f^* \\ & & \text{Psh } Y \end{array}$$

- for $X \in \text{Cat}$, the 2-cell $\theta_X : y_X^* \rightarrow 1$ is also uniquely determined by the universal property of the left Kan extension:

$$\begin{array}{ccc} X & \xrightarrow{y} & \text{Psh } X \\ & \searrow y & \downarrow 1 \\ & & \text{Psh } X \end{array} \quad \begin{array}{ccc} X & \xrightarrow{y} & \text{Psh } X \\ & \searrow y & \downarrow \eta_y \\ & & \text{Psh } X \end{array} \quad \begin{array}{c} \xrightarrow{1} \\ \downarrow 1 \\ \xrightarrow{\theta_X} \\ \downarrow 1 \end{array}$$

The presheaf construction being the free cocompletion, we expect that pseudoalgebras over Psh should be exactly the locally-small cocomplete categories.

Proposition 2.1. *Let $(\mathbb{A}, a; \tilde{a}, \hat{a})$ be a pseudoalgebra for the presheaf relative pseudomonad. Then $\mathbb{A} \in \text{CAT}$ is cocomplete, and the colimit of the diagram $F : \mathbb{D} \rightarrow \mathbb{A}$ for $\mathbb{D} \in \text{Cat}$ is given by*

$$\text{colim } F \cong F^a \text{ colim } Y_{\mathbb{D}}.$$

Proof. Define $s := \text{colim } Y_{\mathbb{D}} \in \text{Psh } \mathbb{D}$ (s standing for ‘singleton’, since it is the terminal presheaf on \mathbb{D} sending every d to a singleton set). Let \mathbb{A} and $F : \mathbb{D} \rightarrow \mathbb{A}$ be as hypothesised. Our proof proceeds as follows:

- (1) Show that $F^a s$ is the apex of a cocone under F .
- (2) Given any cocone under F with apex Gt , construct a map $z_G : F^a s \rightarrow Gt$.
- (3) Show that z_G is a map of cocones.
- (4) Show that if $g : F^a s \rightarrow Gt$ is a map of cocones parallel to z_G , then $g = z_G$.

For (1), consider the colimit cocone under $Y_{\mathbb{D}}$ whose legs $v_d : Yd \rightarrow \text{colim } Y$ are the colimit inclusions. Now for $d \in \mathbb{D}$ we can construct the composite

$$Fd \xrightarrow{\bar{a}} F^a Yd \xrightarrow{F^a v_d} F^a s.$$

Since the following diagram comprising a naturality square and the image under F^a of a cocone:

$$\begin{array}{ccc} Fd & \xrightarrow{Ff} & Fd' \\ \bar{a} \downarrow & & \downarrow \bar{a} \\ F^a Yd & \xrightarrow{F^a Yf} & F^a Yd' \\ & \searrow F^a v_d & \swarrow F^a v_{d'} \\ & F^a s & \end{array}$$

commutes for all $f : d \rightarrow d'$, these form the legs of the required cocone under F with apex $F^a s$.

For (2), we begin by characterising cocones under F . Define $\mathbb{D}^t \in \text{Cat}$ to be the category formed by freely adding a terminal object to \mathbb{D} ; that is,

$$\text{ob } \mathbb{D}^t = \text{ob } \mathbb{D} \sqcup \{t\},$$

$$\text{mor } \mathbb{D}^t = \text{mor } \mathbb{D} \sqcup \{1_t\} \sqcup \{d \xrightarrow{!_d} t : d \in \text{ob } \mathbb{D}\}.$$

We have an inclusion $i : \mathbb{D} \rightarrow \mathbb{D}^t$, and functors $G : \mathbb{D}^t \rightarrow \mathbb{A}$ such that $Gi = F$ correspond exactly to cocones under F .

Let $G : \mathbb{D}^t \rightarrow \mathbb{A}$ be such a functor; we need to construct a map $z_G : F^a s \rightarrow Gt$. Consider the following commutative diagram:

$$\begin{array}{ccc} \mathbb{D} & \xrightarrow{Y} & \text{Psh } \mathbb{D} \\ \downarrow i & \searrow Y_i \xrightarrow{\eta_{Y_i}} & \downarrow (Y_i)^* \\ \mathbb{D}^t & \xrightarrow{Y} & \text{Psh } \mathbb{D}^t \\ & \searrow G \xrightarrow{\bar{a}} & \downarrow G^a \\ & & \mathbb{A} \end{array}$$

and consider the objects $s \in \text{Psh } \mathbb{D}$ and $Yt \in \text{Psh } \mathbb{D}^t$. Since t is terminal in \mathbb{D}^t , Yt is terminal in $\text{Psh } \mathbb{D}^t$. Define the presheaf

$$X := (Y_i)^* s$$

in $\text{Psh } \mathbb{D}^t$ (explicitly, this is the presheaf that sends objects d to singletons and t to the empty set). Then we have a unique map

$$X \xrightarrow{!} Yt$$

in $\text{Psh } \mathbb{D}^t$. Applying the functor G^a to this map, we obtain

$$G^a X \xrightarrow{G^a !} G^a Yt,$$

which we can compose with the following isomorphisms:

$$F^a s = (Gi)^a s \xrightarrow{\bar{a}} (G^a Y_i)^a s \xrightarrow{\hat{a}} G^a (Y_i)^* s \xrightarrow{G^a !} G^a Yt \xrightarrow{\bar{a}^{-1}} Gt$$

giving us a map $F^a s \rightarrow Gt$, and this is how we define $F^a s \xrightarrow{z_G} Gt$.

For (3), we need to show this is a map of cocones, which is to say that the diagram

$$\begin{array}{ccc}
 Fd & \xlongequal{\quad} & Gid \\
 \bar{a} \downarrow & & \downarrow G! \\
 F^a Yd & & \\
 F^a v_d \downarrow & & \downarrow \\
 F^a s & \xrightarrow{z_G} & Gt
 \end{array}$$

commutes. Writing out the definition of z_G , we can fill this diagram with two equalities and two natural transformations:

$$\begin{array}{ccccc}
 Fd & \xlongequal{\quad} & Gid & \xrightarrow{\bar{a}} & G^a Yid \\
 \bar{a} \downarrow & & \bar{a} \downarrow & & \bar{a} \downarrow \\
 F^a Yd & \xlongequal{\quad} & (Gi)^a Yd & \xrightarrow{\bar{a}} & (G^a Yi)^a Yd \\
 F^a v_d \downarrow & & v_d \downarrow & & v_d \downarrow \\
 F^a s & \xlongequal{\quad} & (Gi)^a s & \xrightarrow{\bar{a}} & (G^a Yi)^a s
 \end{array}$$

as well as two naturality squares, pseudoalgebra coherence (5) and the image of a unique map into Yt under G^a :

$$\begin{array}{ccccccc}
 & & G^a Yid & & & & \\
 & & \bar{a} \downarrow & & & & \\
 & & (G^a Yi)^a Yd & \xrightarrow{\bar{a}} & G^a (Yi)^* Yd & \xrightarrow{\eta^{-1}} & G^a Yid & \xrightarrow{\bar{a}^{-1}} & Gid \\
 & & v_d \downarrow & & v_d \downarrow & & \downarrow G^a Y! & & \downarrow G! \\
 & & (G^a Yi)^a s & \xrightarrow{\bar{a}} & G^a (Yi)^* s & \xrightarrow{G^a!} & G^a Yt & \xrightarrow{\bar{a}^{-1}} & Gt
 \end{array}$$

Hence $F^a s \xrightarrow{z_G} Gt$ is a map of cocones.

This gives us that $F^a s$ is a weak colimit for the diagram F ; it remains to do (4) and show uniqueness. Let $G : \mathbb{D}^t \rightarrow \mathbb{A}$ correspond to a cocone under F , and let $g : F^a s \rightarrow Gt$ be a map of cocones; we want to show that $g = z_G$. Define $H : \mathbb{D}^t \rightarrow \text{Psh } \mathbb{D}$ as follows:

$$\begin{aligned}
 H &: \mathbb{D}^t \rightarrow \text{Psh } \mathbb{D} \\
 id &\mapsto Yd \\
 t &\mapsto s \\
 (id \xrightarrow{!} t) &\mapsto (Yd \xrightarrow{v_d} s),
 \end{aligned}$$

noting that $Hi = Y_{\mathbb{D}}$. Then we can define a 2-cell $\beta : F^a H \rightrightarrows G$ with components

$$\begin{aligned}
 \beta_{id} &: F^a Hid = F^a Yd \xrightarrow{\bar{a}^{-1}} Fd = Gid, \\
 \beta_t &: F^a Ht = F^a s \xrightarrow{g} Gt;
 \end{aligned}$$

the only nontrivial naturality condition to check is

$$\begin{array}{ccc} F^a H i d & \xrightarrow{\beta_{i d}} & G i d \\ F^a H ! \downarrow & & \downarrow G ! \\ F^a H t & \xrightarrow{\beta_t} & G t \end{array}$$

and this expands to

$$\begin{array}{ccc} F^a Y d & \xrightarrow{\hat{a}^{-1}} & F d \equiv G i d \\ F^a v_d \downarrow & & \downarrow G ! \\ F^a s & \xrightarrow{g} & G t \end{array}$$

which is precisely the condition that g is a map of cocones. Now, writing out the definition of z_G we obtain

$$F^a s \equiv (G i)^a s \xrightarrow{\tilde{a}} (G^a Y i)^a s \xrightarrow{\hat{a}} G^a (Y i)^* s \xrightarrow{G^a !} G^a Y t \xrightarrow{\tilde{a}^{-1}} G t$$

We use three naturality squares and pseudoalgebra coherence (5):

$$\begin{array}{ccccc} F^a H^*(Y i)^* s & \xrightarrow{F^a H^* !} & F^a H^* Y t & & \\ \hat{a}^{-1} \downarrow & & \hat{a}^{-1} \downarrow & \searrow \eta^{-1} & \\ (F^a H)^a (Y i)^* s & \xrightarrow{(F^a H)^a !} & (F^a H)^a Y t & \xrightarrow{\tilde{a}^{-1}} & F^a H t \\ \beta \downarrow & & \beta \downarrow & & \downarrow \beta_t = g \\ G^a (Y i)^* s & \xrightarrow{G^a !} & G^a Y t & \xrightarrow{\tilde{a}^{-1}} & G t \end{array}$$

along with four naturality squares and pseudoalgebra coherence (3):

$$\begin{array}{ccccc} F^a (H i)^* s & \xrightarrow{\eta} & F^a (H^* Y i)^* s & \xrightarrow{\mu} & F^a H^*(Y i)^* s \\ \hat{a}^{-1} \downarrow & & \downarrow \hat{a}^{-1} & & \downarrow \hat{a}^{-1} \\ (F^a H i)^a s & \xrightarrow{\eta} & (F^a H^* Y i)^a s & & \\ \parallel & & \downarrow \hat{a}^{-1} & & \downarrow \hat{a}^{-1} \\ (F^a H i)^a s & \xrightarrow{\tilde{a}} & ((F^a H)^a Y i)^a s & \xrightarrow{\tilde{a}} & (F^a H)^a (Y i)^* s \\ \beta \downarrow & & \downarrow \beta & & \downarrow \beta \\ (G i)^a s & \xrightarrow{\tilde{a}} & (G^a Y i)^a s & \xrightarrow{\tilde{a}} & G^a (Y i)^* s \end{array}$$

and finally two equalities and pseudoalgebra coherence (4):

$$\begin{array}{ccc} F^a s & \xrightarrow{\theta^{-1}} & F^a Y^* s \equiv F^a (H i)^* s \\ & \searrow & \downarrow \hat{a}^{-1} \quad \downarrow \hat{a}^{-1} \\ & & (F^a Y)^a s \equiv (F^a H i)^a s \\ & & \downarrow \tilde{a}^{-1} \quad \downarrow \beta \\ & & F^a s \equiv (G i)^a s \end{array}$$

which allows us to rewrite the composite defining z_G as:

$$F^a s \xrightarrow{\theta^{-1}} F^a Y^* s = F^a (H i)^* s \xrightarrow{\eta} F^a (H^* Y i)^* s \xrightarrow{\mu} F^a H^*(Y i)^* s \xrightarrow{F^a H^* Y !} F^a H^* Y t \xrightarrow{\eta^{-1}} F^a H t \xrightarrow{g} G t.$$

Now this is of the form

$$F^a s \xrightarrow{F^a(\dots)} F^a s \xrightarrow{g} Gt$$

for some map $s \rightarrow s$. But since s is terminal in $\text{Psh } \mathbb{D}$, the only such map is 1_s . Hence by functoriality we have

$$z_G = g \circ F^a(1_s) = g \circ 1_{F^a s} = g.$$

So indeed the map of cocones $F^a s \rightarrow Gt$ is unique, which implies

$$F^a s = F^a \text{colim } Y_{\mathbb{D}} \cong \text{colim } F.$$

Hence every presheaf pseudoalgebra \mathbb{A} is cocomplete. □

REFERENCES

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